

Optimizing till stationarity on the bounded-rank variety

Nicolas Boumal

EPFL, Institute of Mathematics

Ongoing work with Eitan Levin (CalTech) and Joe Kileel (UT Austin)

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My focus today: bounded rank constraints

$$\min_{X \in \mathbf{R}^{m \times n}} f(X) \quad \text{subject to} \quad \text{rank}(X) \leq k$$

Many applications (some with additional structure, e.g., $X \succcurlyeq 0$):

Model order reduction

Recommender systems

Sensor network localization

Large scale linear matrix equations

...

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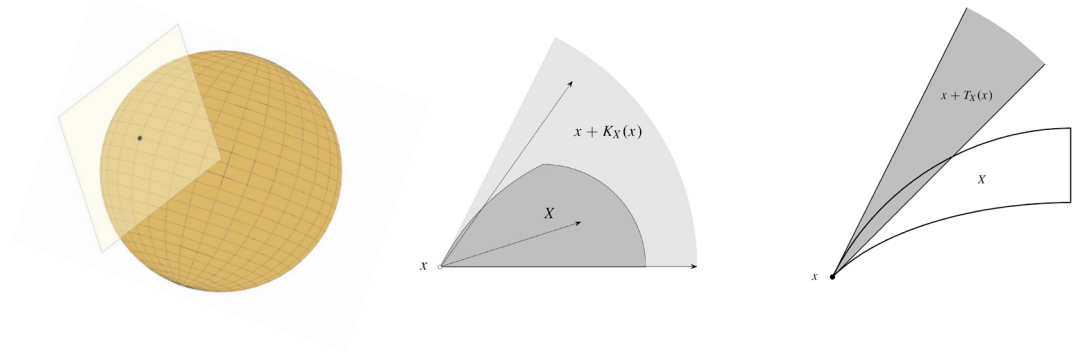
Assume f is smooth (let's say C^∞).

In general, finding a global minimum is NP-hard.*

Less ambitious goal of this talk: **find a stationary point.**

*See for example Gillis & Glineur, *Low-rank matrix approximation with weights or missing data is NP-hard*, SIMAX 2011

Stationarity in general



$$\min_{x \in \mathcal{X}} f(x) \quad \text{subject to } x \in \mathcal{X}$$

The **tangent cone** $T_x \mathcal{X}$ collects allowed directions of movement at x .

$$T_x \mathcal{X} = \left\{ \lim_{i \rightarrow \infty} \frac{x_i - x}{\tau_i} : (x_i) \subset \mathcal{X}, \tau_i \in \mathbf{R}^+, x_i \rightarrow x, \tau_i \rightarrow 0 \right\}$$

x is **stationary** if $Df(x)[v] \geq 0$ for all $v \in T_x \mathcal{X}$, i.e., $-\nabla f(x) \in (T_x \mathcal{X})^\circ$.

This is equivalent to the property $\|\text{Proj}_{T_x \mathcal{X}}(-\nabla f(x))\| = 0$.

Rank **bound** or **equality**: different geometries

The following set **is a smooth manifold**:

$$\{X \in \mathbf{R}^{m \times n} : \text{rank}(X) = k\}$$

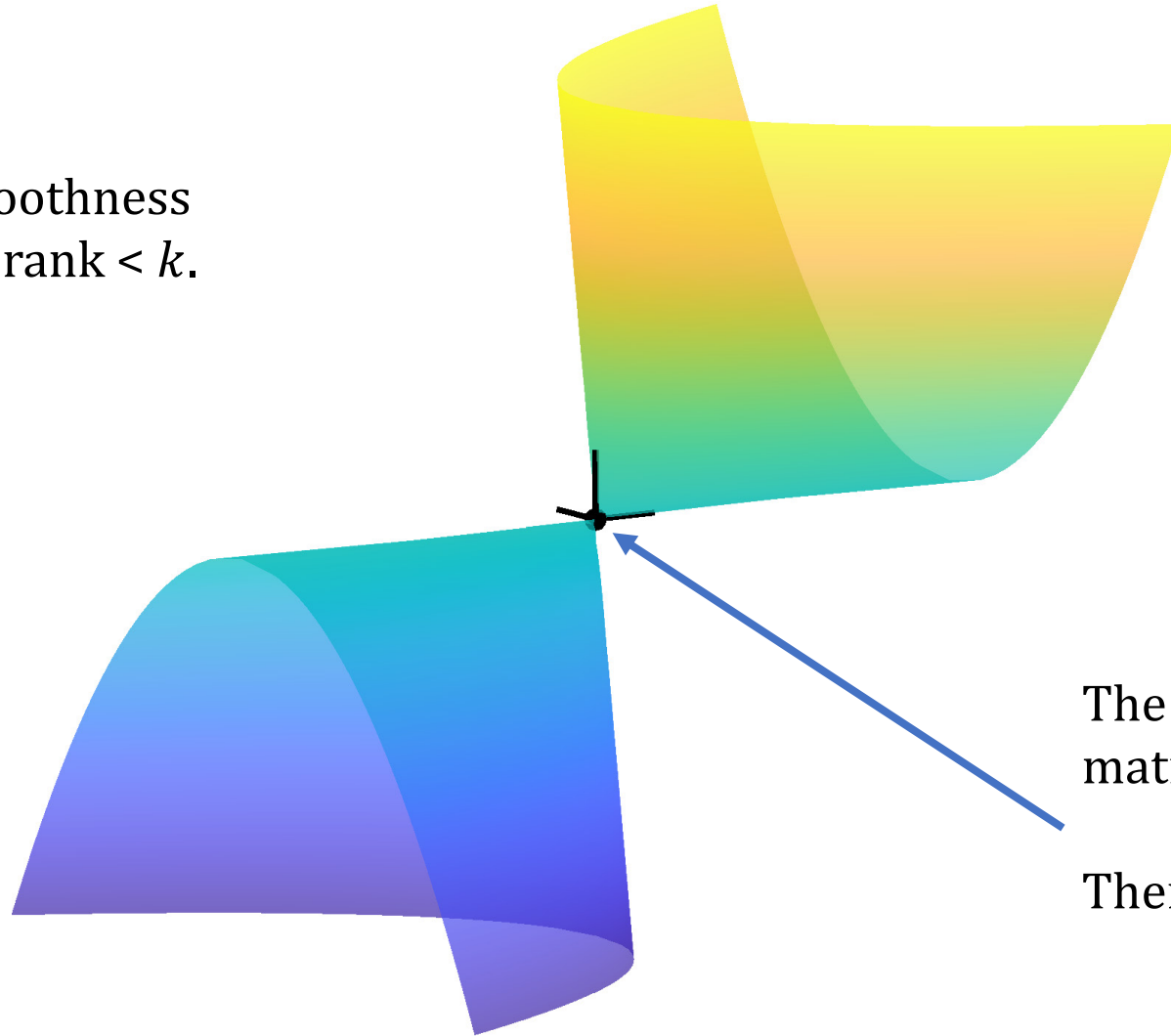
However, the following set **is not**:

$$\{X \in \mathbf{R}^{m \times n} : \text{rank}(X) \leq k\}$$

Let's do a **proof by picture** for related case of symmetric $X \in \mathbf{R}^{2 \times 2}$.

$$\{X \in \mathbf{R}^{2 \times 2} : X = X^T \text{ and } \text{rank}(X) \leq 1\} = \left\{ \begin{bmatrix} x & y \\ y & z \end{bmatrix} : xz - y^2 = 0 \right\}$$

More generally, smoothness fails at all points of rank $< k$.



The origin is the only matrix of rank zero.

There, the set is not smooth.

Rank constraints: mind the cliff

$$\min_{X \in \mathbb{R}^{m \times n}} f(X) \quad \text{subject to} \quad \text{rank}(X) \leq k$$

If the iterates remain comfortably on the manifold of rank- k , fine. In practice, this is often the case.

However, if the iterates approach lesser-rank matrices, or worse, converge to one, then smooth optimization theory breaks down.

Optimization algorithms must be able to handle this eventuality.

Even computing stationary points is tricky

$$\min_{X \in \mathbf{R}^{m \times n}} f(X) \quad \text{subject to} \quad \text{rank}(X) \leq k$$

projected gradient descent method*

$$X_{i+1} = \text{Proj}_{\mathbf{R}_{\leq k}^{m \times n}} \left(X_i + \alpha_i \text{Proj}_{T_{X_i} \mathbf{R}_{\leq k}^{m \times n}} (-\nabla f(X_i)) \right)$$

*Schneider & Uschmajew, SIOPT 2015,

Convergence Results for Projected Line-Search Methods on Varieties of Low-Rank Matrices Via Łojasiewicz Inequality

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There exist f and X_0 for which a **projected gradient descent** method* with Armijo backtracking produces iterates X_1, X_2, X_3, \dots such that:

1. $\text{rank}(X_i) = k$ for all i ,
2. The **stationarity measure goes to zero** as $i \rightarrow \infty$,
3. The sequence converges to a feasible matrix X ,
4. Yet **the limit X is not stationary**. – We might be far from any!

*Schneider & Uschmajew, SIOPT 2015,

Convergence Results for Projected Line-Search Methods on Varieties of Low-Rank Matrices Via Łojasiewicz Inequality

Apocalypses in general (algorithm agnostic)

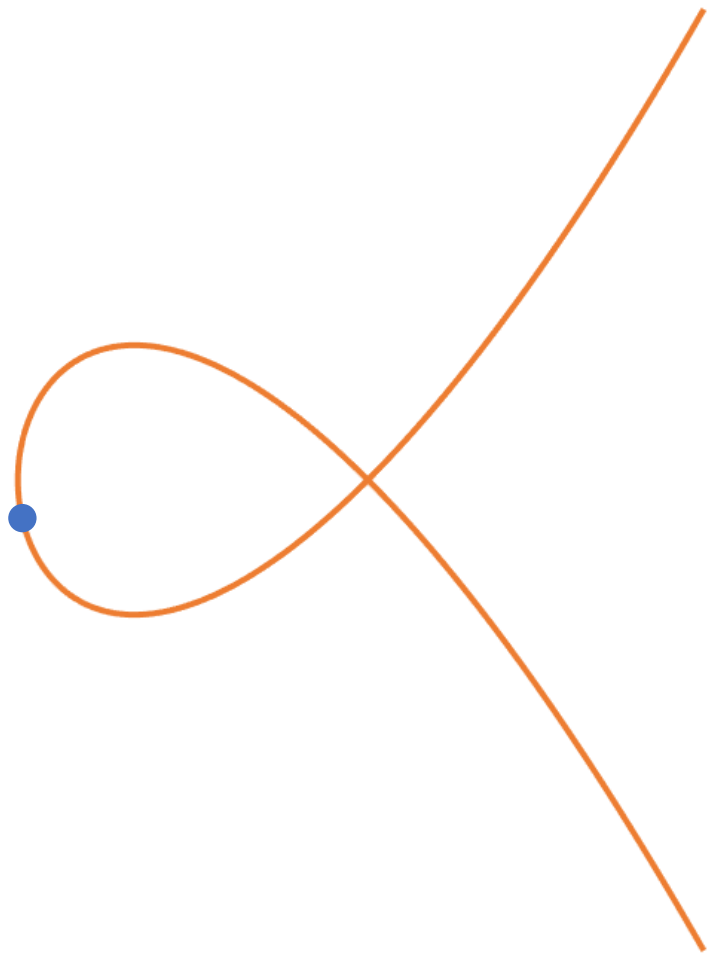
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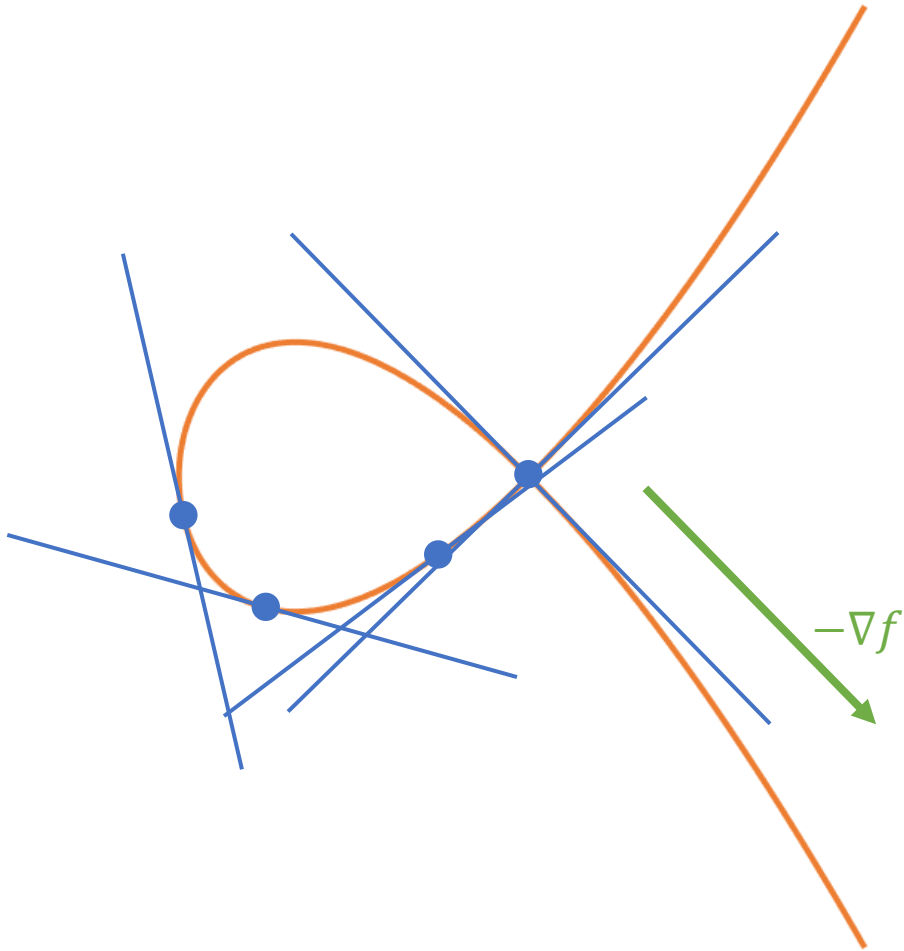
x is **apocalyptic** if there exists a sequence $x_i \rightarrow x$ and a function f such that $\|\text{Proj}_{T_{x_i} \mathcal{X}}(-\nabla f(x_i))\| \rightarrow 0$ yet x is not stationary.

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When can apocalypses occur?

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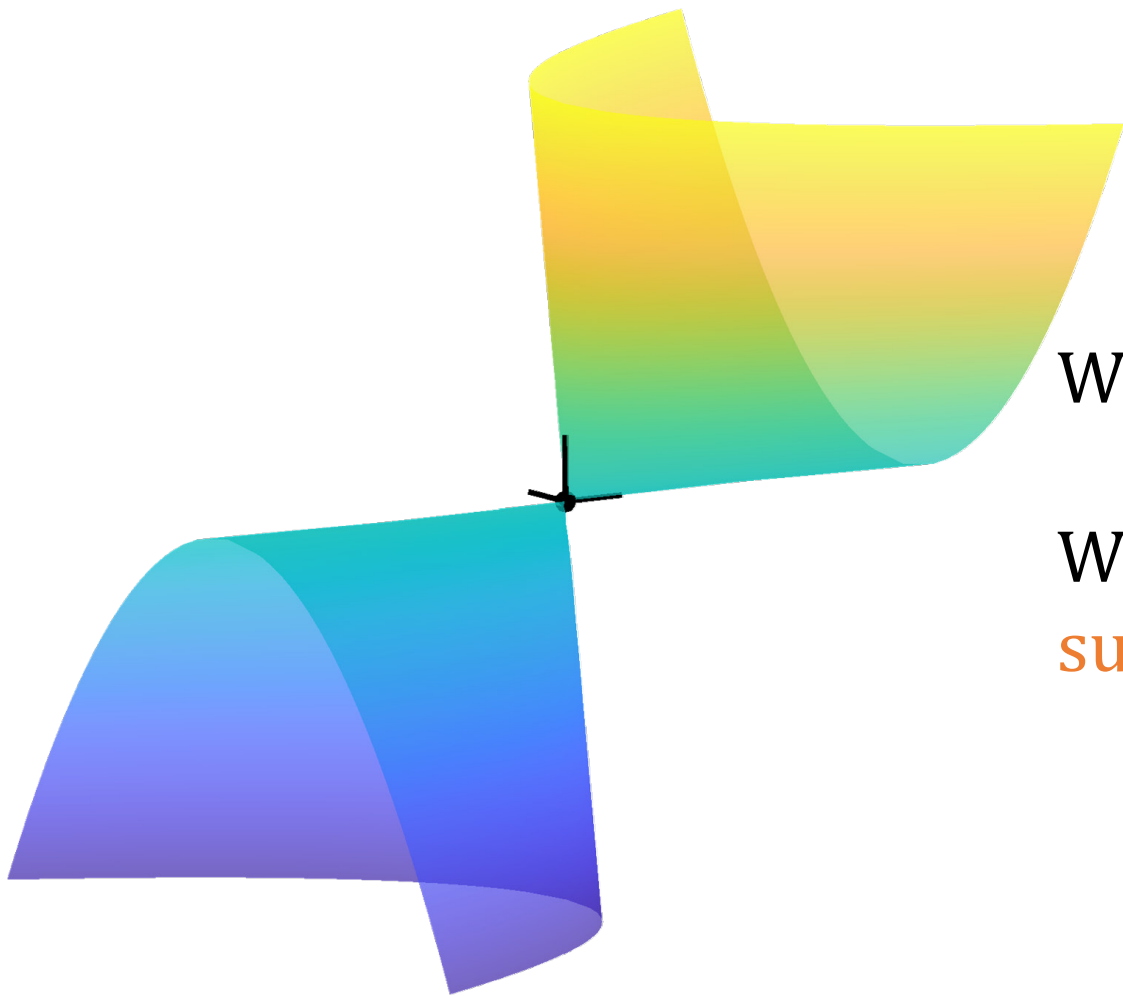


When can apocalypses occur?

When tangent cones change
suddenly, adding **new directions**:

$$\left(\limsup_{i \rightarrow \infty} T_{x_i} \mathcal{X} \right)^\circ \not\subseteq (T_x \mathcal{X})^\circ$$

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A first take-away, and two positive notes

Apocalypses are a geometric feature of a search space \mathcal{X} .

They **cause blind spots** for gradient-based methods.

Two positive notes:

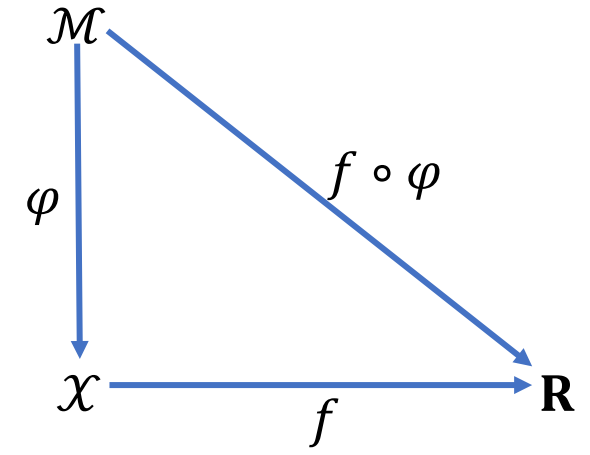
Convex sets and **manifolds with boundaries** have no apocalypses.

Using **second-order** information of f , we can find stationary points.

To find stationary points, use lifts

Let $\mathcal{M} = \mathbf{R}^{m \times k} \times \mathbf{R}^{n \times k}$ and $\mathcal{E} = \mathbf{R}^{m \times n}$.

Consider the smooth map $\varphi(L, R) = LR^\top$ from \mathcal{M} to \mathcal{E} .



Notice: $\varphi(\mathcal{M}) = \{X \in \mathbf{R}^{m \times n} : \text{rank}(X) \leq k\}$: it is a **smooth lift**.

Thm*: If (L, R) is 2-critical for $f \circ \varphi$, then LR^\top is stationary for f .

Thm: If f has compact sublevel sets, then a modified version of the trust-region method on $f \circ \varphi$ finds 2-critical points, always.

* Ha, Liu & Barber, SIOPT 2021,

An equivalence between critical points for rank constraints versus low-rank factorizations

Summary

Optimization on nonsmooth sets: watch out for apocalypses.

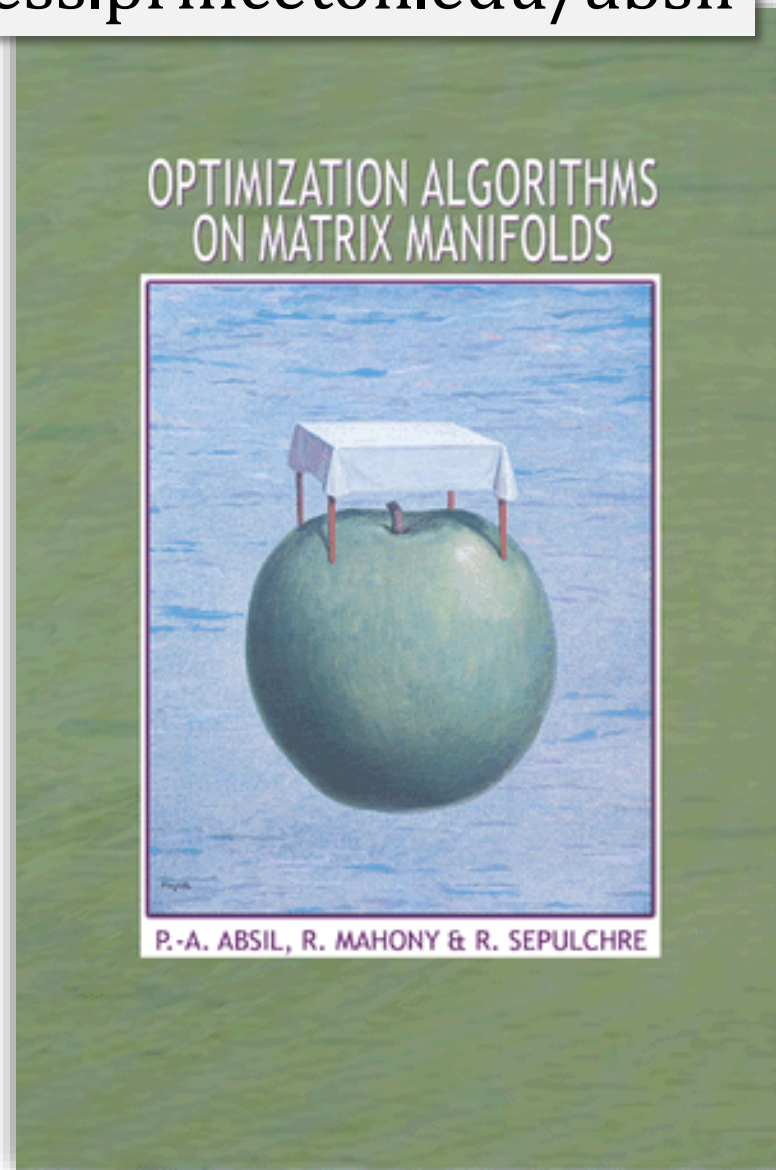
Exist on bounded-rank variety; not on convex sets / manifolds with boundary.

We can use lifts to move the problem to a smooth manifold.

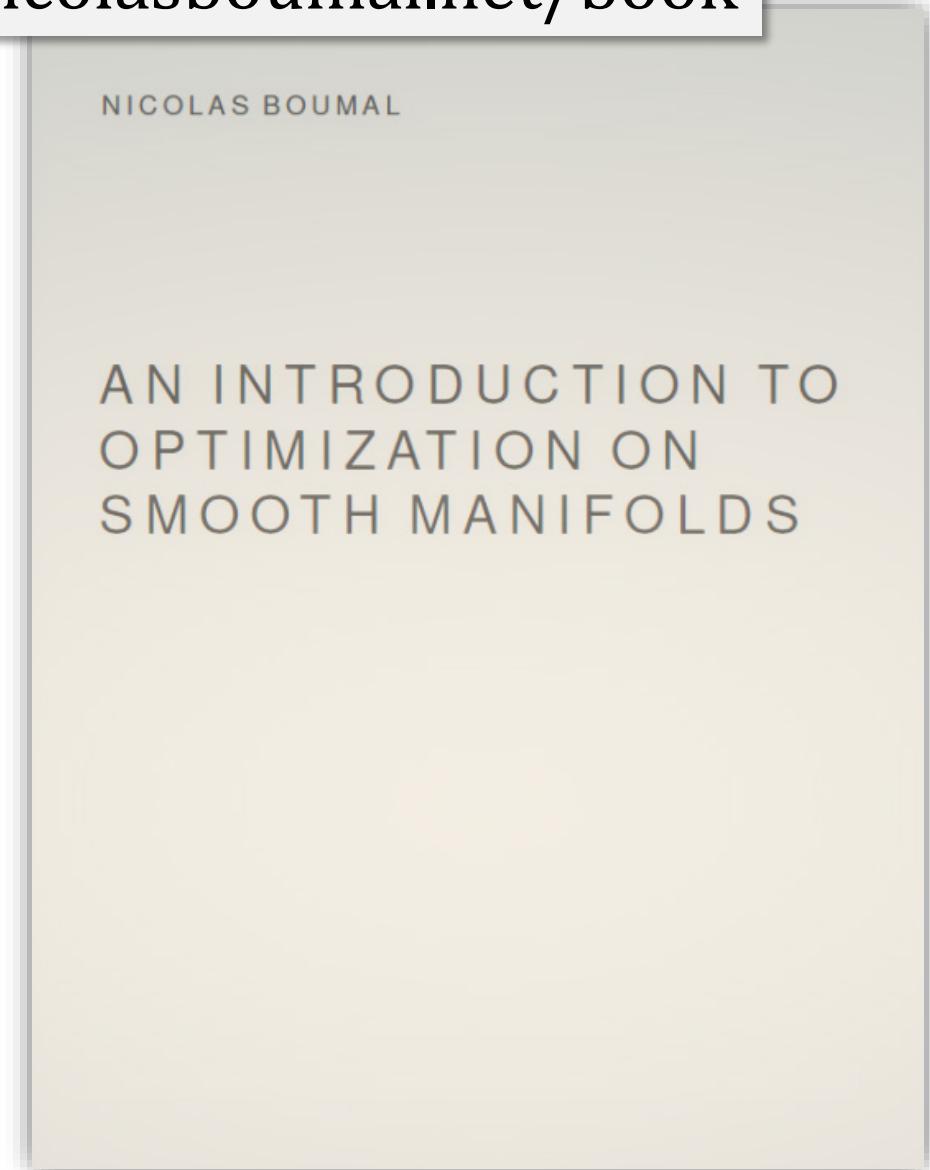
Can be done for other nonsmooth sets: desingularization, symmetry, shadows...

If the lift has nice properties (e.g., 2-critical \mapsto stationary), it tells us how to use the Hessian of f to avoid apocalypses.

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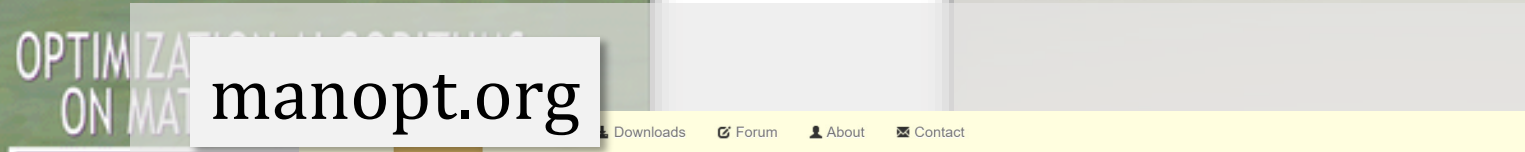
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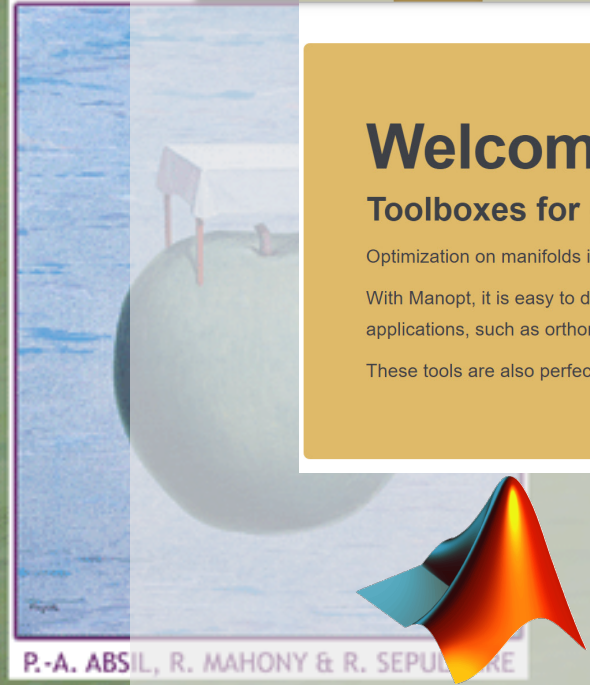
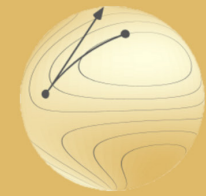
Welcome to Manopt!

Toolboxes for optimization on manifolds and matrices

Optimization on manifolds is a powerful paradigm to address nonlinear optimization problems.

With Manopt, it is easy to deal with various types of constraints and symmetries which arise naturally in applications, such as orthonormality, low rank, positivity and invariance under group actions.

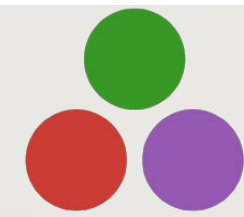
These tools are also perfectly suited for unconstrained optimization with vectors and matrices.



With Bamdev Mishra,
P.-A. Absil & R. Sepulchre



Lead by J. Townsend,
N. Koep & S. Weichwald



Lead by Ronny Bergmann