

Low-rank approaches for SDP's in community detection



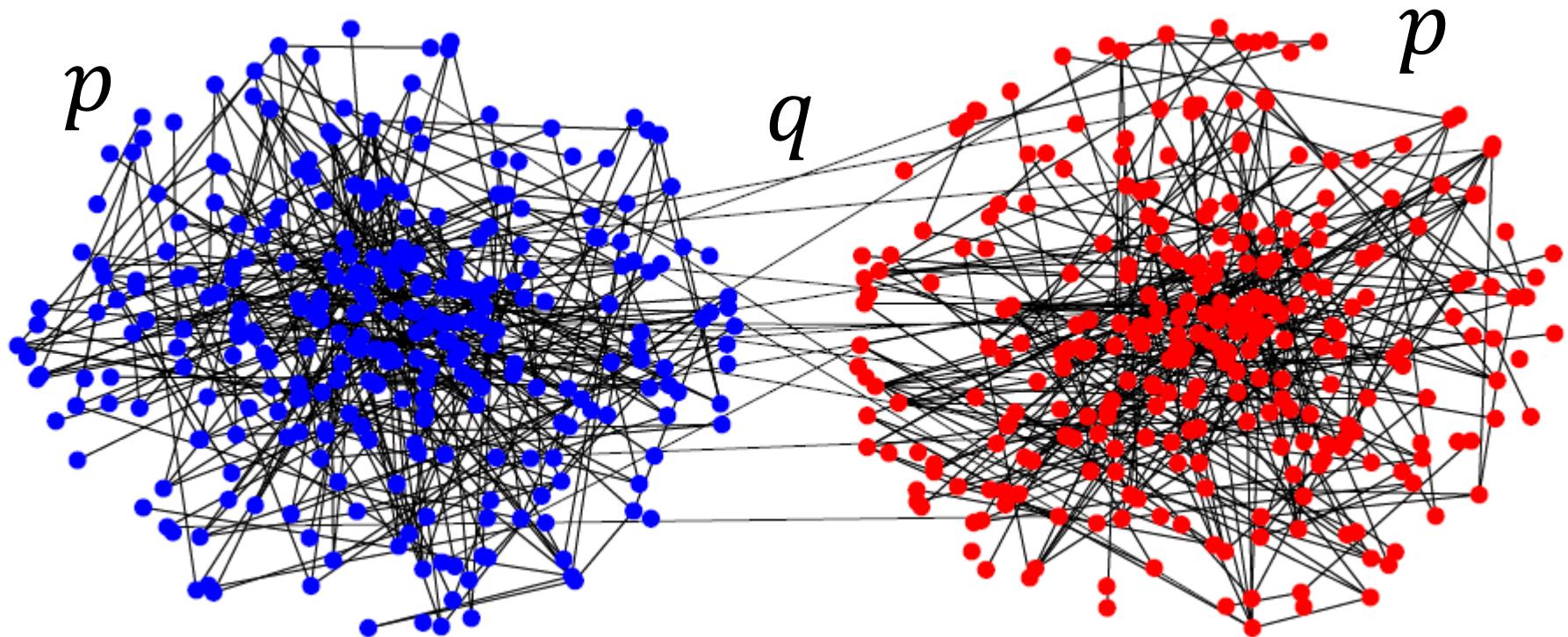
Afonso Bandeira (MIT→Courant)

Nicolas Boumal (Princeton University)

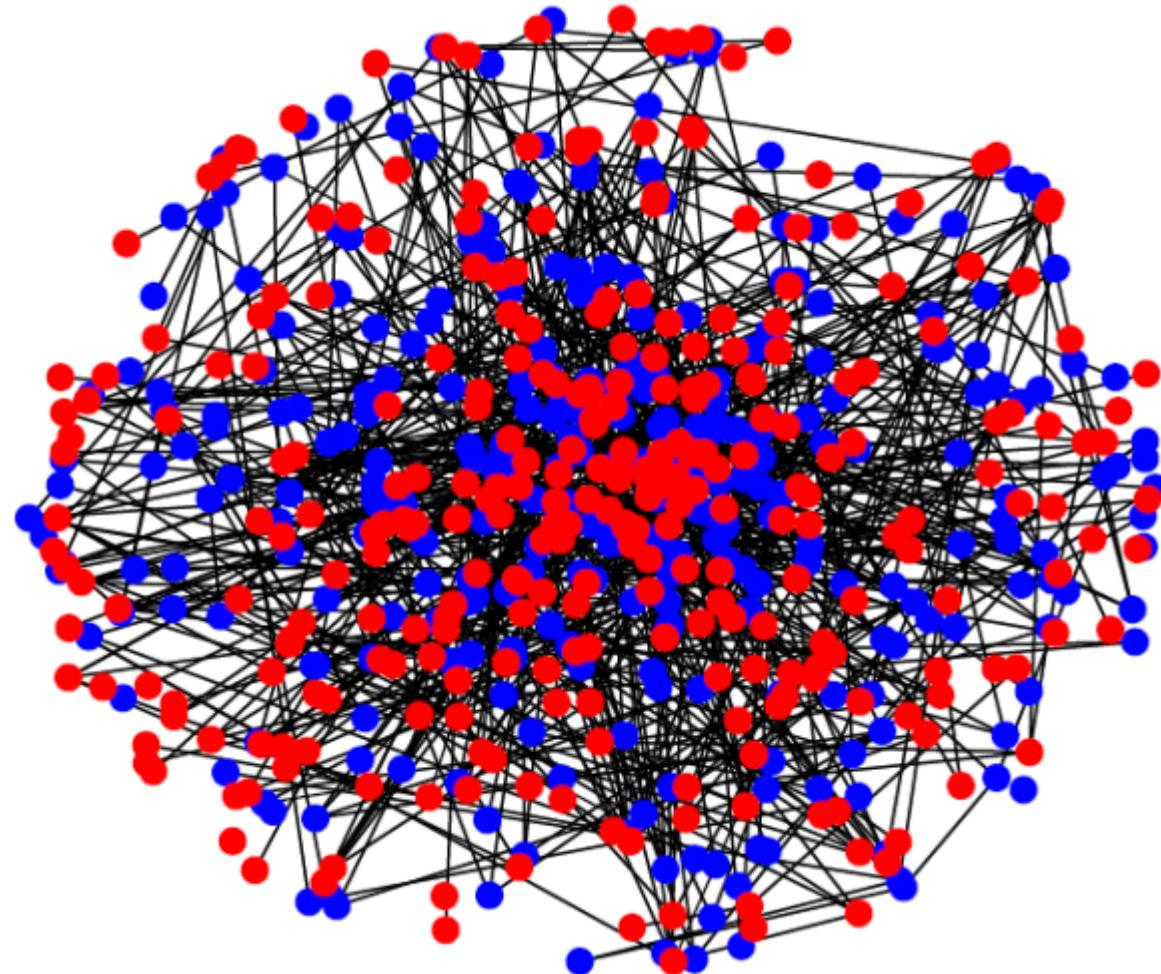


Vladislav Voroninski (MIT)

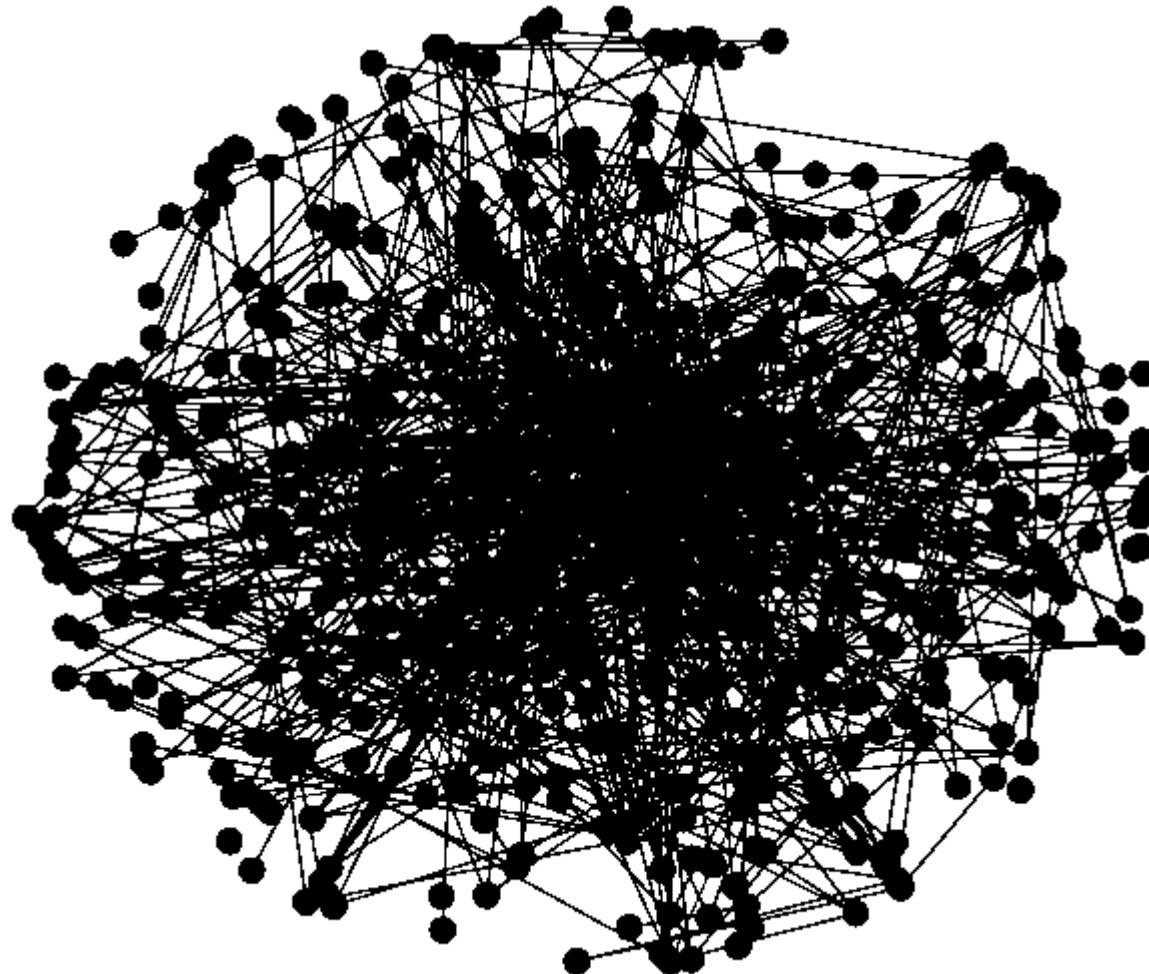
Community detection in the stochastic block model (SBM)



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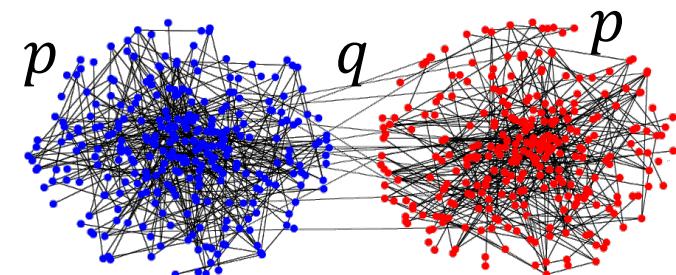
The constant average degree regime

Link **within**: $p = \frac{a}{n}$

Link **across**: $q = \frac{b}{n}$

Each community has a giant connected component, whp.

Hope for **non-trivial correlation** with true partition.



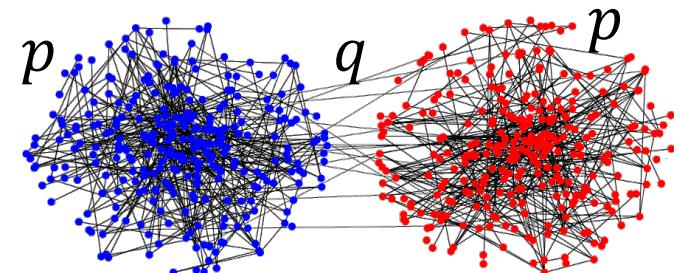
The dense regime

Link **within**: $p = \alpha \frac{\log n}{n}$

Link **across**: $q = \beta \frac{\log n}{n}$

Each community is connected, whp.

Hope for **exact recovery**.



The key SNR quantity: $\lambda(p, q)$

$$\lambda(p, q) = \frac{p - q}{\sqrt{2(p + q)}} \sqrt{n}$$

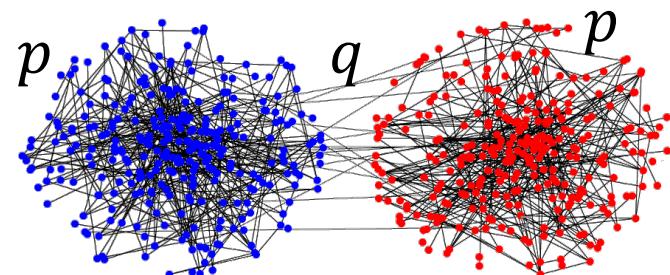
For non-trivial correlation: need $\lambda > 1$

Decelle et al. '11, Mossel et al. '14, Massoulié et al. '14

For exact recovery: “need $\lambda > \sqrt{2 \log n}$ ”

Mossel et al. '14, Abbe et al. '14

(Precise condition: $\sqrt{\alpha} - \sqrt{\beta} > \sqrt{2}$.)



Relaxation of MLE gives SDP for SBM

With A the adjacency matrix and $\textcolor{brown}{A}' = A - \frac{p+q}{2} \mathbf{1}\mathbf{1}^T$:

$$\max_X \langle \textcolor{brown}{A}', X \rangle \text{ s.t. } \text{diag}(X) = \mathbf{1}, X \geq 0$$

Non-trivial correlation:

$\forall \delta > 0$, if $\frac{p+q}{2}n$ large enough, $\lambda > 1 + \delta$ is enough

Guedon & Vershynin '14, Montanari & Sen '15, Javanmard et al. '15

Exact recovery: SDP is tight **at the info limit**

Hajek et al. '14, Bandeira '15

The Burer–Monteiro approach

$$\max_X \langle A', X \rangle \text{ s.t. } \text{diag}(X) = \mathbf{1}, X \geq 0$$

Parameterize $X = YY^T$ with Y of size $n \times p$:

$$\max_Y \langle A', YY^T \rangle \text{ s.t. } \text{diag}(YY^T) = \mathbf{1}$$

The aggressive version: $p = 2$.

Non-convex optimization on the n -torus

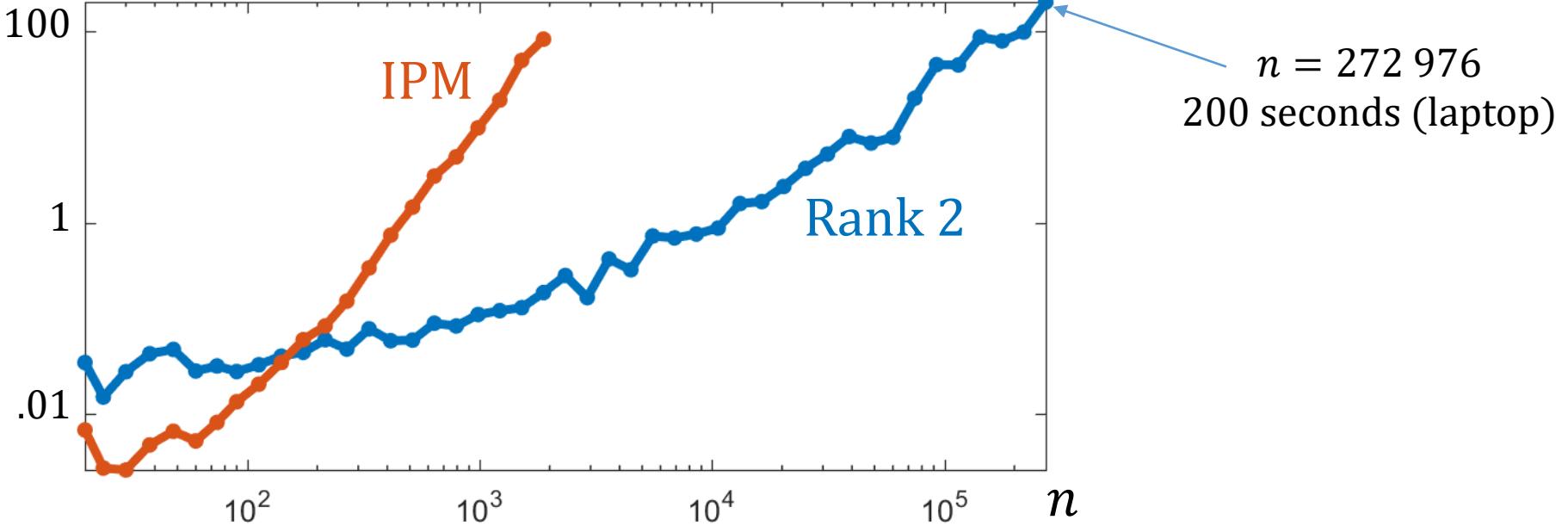
$$\max_{Y \in \mathbb{R}^{n \times 2}} \langle A', YY^T \rangle \text{ s.t. } \text{diag}(YY^T) = \mathbf{1}$$

Low-dimensional, and no conic constraint.

We run Riemannian trust regions via Manopt.

Do KKT points have good statistical properties?

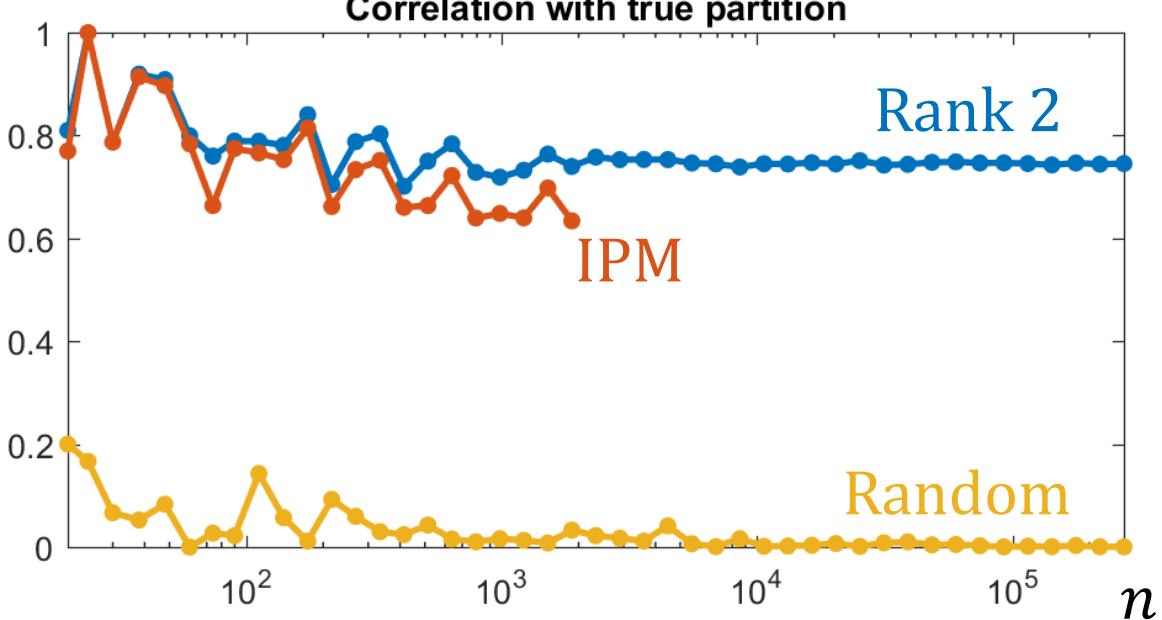
Computation time [s]



$$p = \frac{10}{n}, q = \frac{2}{n}$$

$$\lambda = 1.63$$

Correlation with true partition



Main result 1: non-trivial correlation

$$\max_{Y \in \mathbf{R}^{n \times 2}} \langle A', YY^T \rangle \text{ s.t. } \text{diag}(YY^T) = \mathbf{1}$$

In the constant average degree regime, for any $\delta > 0$, if $\frac{p+q}{2}n$ is large enough and $\lambda > 8 + \delta$,

Then, there exists $\varepsilon > 0$ such that, whp, all second order KKT points Y correlate non-trivially with the true partition \mathbf{g} :

$$\frac{1}{n} \|Y^T \mathbf{g}\|_2 \geq \varepsilon.$$

Main result 2: exact recovery

$$\max_{Y \in \mathbb{R}^{n \times 2}} \langle A', YY^T \rangle \text{ s.t. } \text{diag}(YY^T) = \mathbf{1}$$

There exists c (universal) such that, if

$$\lambda \geq cn^{1/3},$$

then, whp, all second-order KKT points Y are optimal and correspond to \mathbf{g} :

$$YY^T = \mathbf{g}\mathbf{g}^T.$$

$$1. \text{ Hess } f(Y) \geq 0 \Leftrightarrow \text{ddiag}(A'YY^T) \geq A' \circ YY^T$$

$$\langle \text{ddiag}(A'YY^T), \mathbf{g} \mathbf{g}^T \circ YY^T \rangle \geq \langle A' \circ YY^T, \mathbf{g} \mathbf{g}^T \circ YY^T \rangle$$

$$2. \text{ Link } A' \text{ to the signal: } A' \propto \mathbf{g} \mathbf{g}^T + \frac{n}{\lambda} E + D$$

$$\left\langle \mathbf{g} \mathbf{g}^T + \frac{n}{\lambda} E, YY^T \right\rangle \geq \left\langle \mathbf{g} \mathbf{g}^T + \frac{n}{\lambda} E, \mathbf{g} \mathbf{g}^T \circ YY^T \circ YY^T \right\rangle$$

$$3. \text{ Use noise property: } \max_{X \succcurlyeq 0, \text{diag}(X)=\mathbf{1}} \langle E, X \rangle \leq (2 + o_d(1))n$$

$$\|Y^T \mathbf{g}\|^2 \geq \|YY^T\|_{\text{F}}^2 - \frac{2n^2}{\lambda} (2 + o_d(1)) \geq n^2 \left(\frac{1}{2} - \frac{4 + o_d(1)}{\lambda} \right)$$

Note: Sufficient to ensure $\text{Hess } f(Y) \geq -\varepsilon \cdot \mathbf{I}$

More of this in statistics?

It's a tempting family of estimators.

Toward computation bounds: see global rates of convergence to KKT points on manifolds: 1605.08101

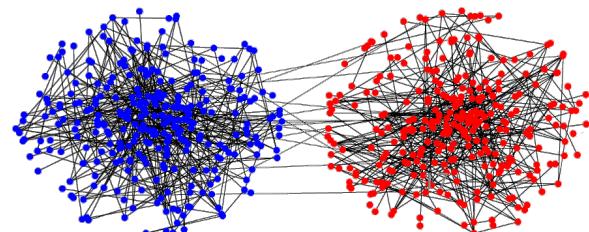
More guarantees for Burer–Monteiro approach to ‘smooth’ SDP’s: 1606.04970

Maximum likelihood estimation for SBM
is combinatorial optimization

$A \in \mathbf{R}^{n \times n}$ is the observed adjacency matrix:

$$\max_{\mathbf{z}} \langle A, \mathbf{z}\mathbf{z}^T \rangle \text{ s.t. } \mathbf{z} \in \{\pm 1\}^n \text{ and } \mathbf{1}^T \mathbf{z} = 0$$

Hard problem for general A : need to relax.



Step 1: remove the linear constraint

$$\max_{\mathbf{z}} \langle A, \mathbf{z}\mathbf{z}^T \rangle \text{ s.t. } \mathbf{z} \in \{\pm 1\}^n \text{ and } \mathbf{1}^T \mathbf{z} = 0$$

If $\mathbf{g} \in \{\pm 1\}^n$ is the true partition, then

$$\mathbf{E}\{A\} = \frac{p+q}{2} \mathbf{1}\mathbf{1}^T + \frac{p-q}{2} \mathbf{g}\mathbf{g}^T.$$

Remove the bias toward 1: $A' = A - \frac{p+q}{2} \mathbf{1}\mathbf{1}^T$

Step 2: relax to a semidefinite program

$$\max_{\mathbf{z}} \langle A', \mathbf{z}\mathbf{z}^T \rangle \text{ s.t. } \mathbf{z} \in \{\pm 1\}^n$$

Binary constraints $\mathbf{z} \in \{\pm 1\}^n$
are equivalent to $\text{diag}(\mathbf{z}\mathbf{z}^T) = \mathbf{1}$.

Introduce $X = \mathbf{z}\mathbf{z}^T$. Equivalent problem:

$$\max_X \langle A', X \rangle \text{ s.t. } \text{diag}(X) = \mathbf{1}, X \succeq 0, \text{rank}(X) = 1$$